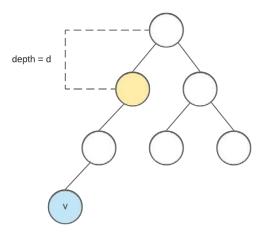
An Analysis of the Level Ancestor Problem

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The Level Ancestor Problem

- In a tree T with n nodes, answer a query of the form LA(v, d)
 - Query aims to find the ancestor of the node *v* that is at a depth of *d* from the root
 - This can be done is O(n) time by tracing the path from the target node to the root until reaching depth *d*, but it is impractical for larger trees
- Preprocessing for the algorithms aims to store the tree in data structures so that queries can be completed more efficiently



Applications

- > Level ancestor implementations can be used to aid in many other algorithms
 - Least Common Ancestor
 - Used in string processing, computational biology, and complex distributed systems, among many other applications

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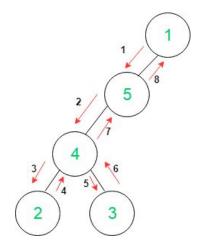
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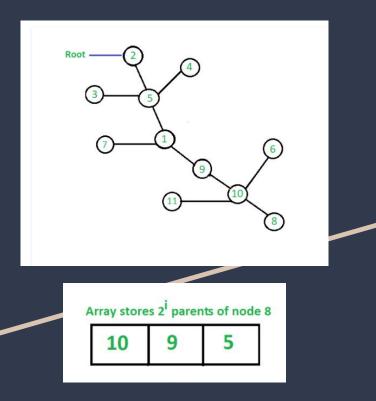
- Space-Efficient Ordinal Trees
 - XML document representation for XPath queries
- Range-Aggregate Queries on trees

Our Project

- We aimed to code (in C) the algorithms outlined in different papers and compare the experimental time for tree preprocessing and querying
- > We examined eight algorithms for static binary trees
- Random Tree Generation:
 - Trees were represented in Euler format that follows a Depth-First Search preorder
 - Binary Representation:
 - 1 represents moving down in the tree
 - 0 represents moving upwards in the tree
 - Example: 11101000
 - Euler Representation= 154243451
 - Generated via split-tree method
 - Creates trees where both the average node depth and tree depth are logarithmically bound
 - Allows for easy construction of Euler representation and skewed trees



The Algorithms



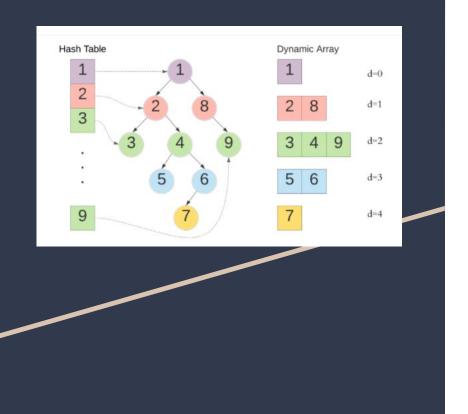
> Table Algorithm:

- Preprocessing = $O(n^2)$, Query = O(1)
- Uses a 2D-array to store the list of all ancestors of each node derived by a traversal in DFS order
- ➢ Jump-Pointer Algorithm:
 - Preprocessing = O(nlogn), Query = O(logn)
 - Uses pointers from each node to its ancestors, allowing for the distance to the ancestor at depth *d* to be traversed in jumps of at least half the remaining distance to the ancestor

➤ Ladder Algorithm:

- Preprocessing = O(n), Query = O(logn)
- Perform a longest path decomposition on the tree and extends the paths to reach the root so that a node at height *h* can be queried for ancestors of at least height 2*h*
- ➤ Jump Ladder:
 - Preprocessing = O(nlogn), Query = O(1)
 - Combines the Jump-Pointer and Ladder Algorithms
 - First jumps up the tree, then uses the ladder to move up towards the ancestor node

The Algorithms



- Macro-Micro Algorithm:
 - Preprocessing = O(n), Query = O(1)
 - Uses the other four algorithms
 - Nodes are divided into macro and micro subtrees based on their weight
 - Macro nodes have use of the jump-pointer method
 - Micro trees use table algorithm
- Menghani-Matani (Google-Facebook) Algorithm:
 - Preprocessing = O(n), Query = O(logn)
 - Perform a pre-order traversal of the tree where each node is given a integer tag starting at 1 for the root
 - Each node, in increasing tag value, is stored in a specific dynamic array depending on its depth
 - Find the largest node label less than the label of node *v* in the array for depth *d* and return the pointer for that label

The Algorithms



- ➤ Torben Hagerup:
 - Preprocessing = O(n), Query = O(1)
 - Uses the find-smaller method to create a 2D structure of peaks and valleys
 - Finds the deepest valley nearest to the node v and then uses the ladder method until reaching the height of the given depth d
 - Then the node to the right of this height in the 2D structure is selected

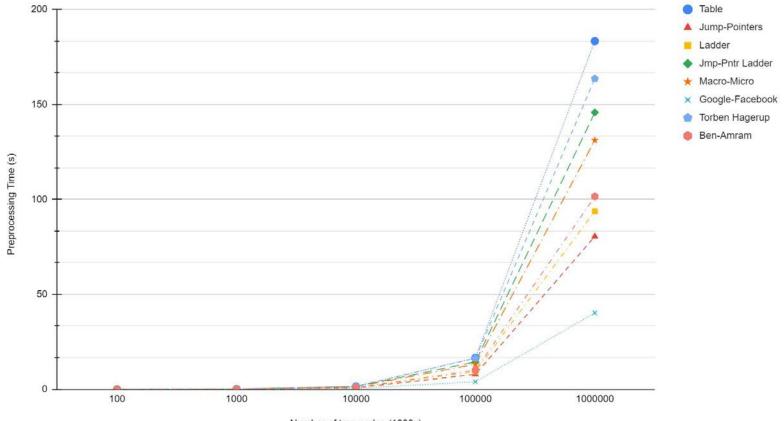
➢ Ben-Amran:

- Preprocessing = O(nlogn), Query = O(n)
- Uses the find-smaller methods which takes an array *A*, an index *x*, and an integer *y*
- Finds an index u > x such that $A_{ij} \le y$

Results

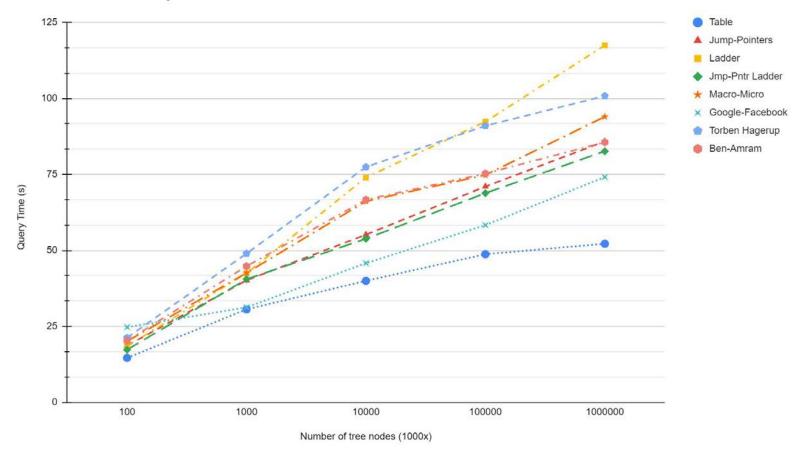
- Multiple trials were run with varying numbers of nodes from 100,000 to 1 billion and 100 million queries per tree
- Tests were run on the TCNJ cluster as submitted job scripts to an AMD EPYC2 processor with 512 GB of memory
- > Preprocessing:
 - Fastest Google-Facebook Algorithm
 - Slowest Table Algorithm
- > Querying:
 - Fastest Table Algorithm
 - Slowest Torben Hagerup & Ladder Algorithms

Level Ancestor Preprocessing Times



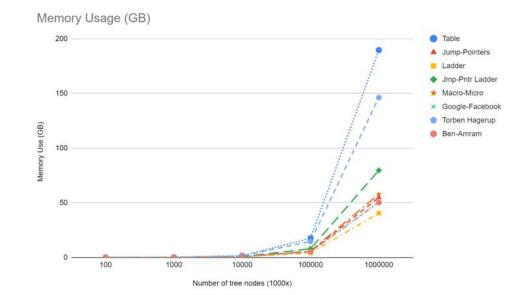
Number of tree nodes (1000x)

Level Ancestor Query Times



Results

- > In addition to time taken, we also examined memory usage for each algorithm
 - As expected, the Table Algorithm uses up the most memory space



Future Research

- Examine ways to improve upon the existing algorithms
 - Could incorporate more assembly code that speeds up processing time
- Incorporate dynamic algorithms that allow for adding leaves to trees
 - Alstrup & Holm Algorithm
 - Implements an improved version of the macro-micro algorithm
 - Dietz Algorithm
 - Adjust static algorithms to allow for dynamic trees

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